Programming Languages (4)
Memory Management Introduction

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Memory management in programming languages

- all values (integers, floating point numbers, strings, arrays, structs, ...) need memory to hold them
- ideally, programming languages manage them on behalf of the programmer
- three approaches covered

manual		C, C++
garbage collection	traversing reference counting	Python, Java, Julia, Go, OCaml, etc.
Rust ownership		Rust

Memory Management Quiz

take the quiz via any of the following

- direct link
- go menti.com and enter code 4574 1905
- use this QR Code

Illustration (Q2)

```
def foo():
    m = node("Mimura")
    o = node("Ohtake")
    o.friend = m
    return o
```

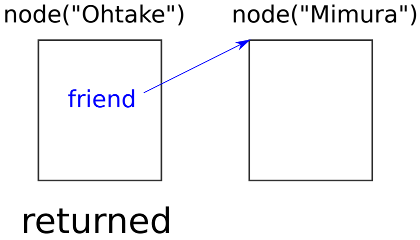

Illustration (Q3)

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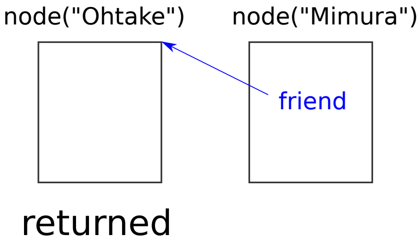

Contents

Memory allocation in C/C++

1. Global variables/arrays
2. Local variables/arrays
3. Heap

```
int g; int ga[10];
int foo() {
  int l; int la[10];
  int * a = &g;
  int * b = ga;
  int * c = &l;
  int * d = la;
  int * e = malloc(sizeof(int));
}
```

- lifetime

	starts	ends
global	when the program starts	when program ends
local	when a block starts	when a block ends
heap	malloc, new	free, delete

- note: the following discussion calls all of them *objects*

What could go wrong in manual memory management (e.g., C/C++)?

- heap-allocated (i.e., `new`/`malloc`'ed) memory must be `delete`/`free`d at the right spot
  - *premature free* = using it after `delete/free` → memory corruption
  - *memory leak* = not `delete/free`ing no-longer-used memory → (eventually) out of memory

```
node * foo() {
  node * m = new node("Mimura");
  node * o = new node("Ohtake");
  return o;
}
```

What could go wrong in manual memory management (e.g., C/C++)?

- stack-allocated memory are automatically reclaimed when it goes out of scope
  - using it afterwards ≡ premature delete

```
node * foo() {
  node m = node("Mimura");
  node o = node("Ohtake");
  return &o;
}
```

```
node * foo() {
  node   m =     node("Mimura");
  node * o = new node("Ohtake");
  o->frien = &m;
  return o;
}
```

Tools to make C/C++ memory management safer

- `valgrind` (memory checker)
  - detect memory-related errors (use after free, memory leak, out of bound accesses, etc.)
- Boehm garbage collection library for C/C++
  - automatically garbage-collect memory blocks allocated by malloc/new

Contents

Garbage Collection (GC)

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  - you may access an object after its lifetime
  - you may not free an object despite you no longer access it

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  - the system automatically does that
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- the question: how does the system know *which objects may be accessed in future*?

Objects that may {ever/never} be accessed

- the precise judgment is undecidable
- (at the start of line 2) "the object pointed to by `p` will ever be accessed" $\iff$ "`f(x)` will terminate and return 0" $\rightarrow$ you need to be able to solve the halting problem...

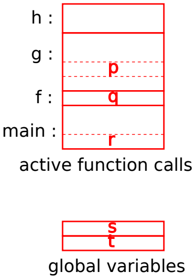
```
1 int main() {
2   if (f(x) == 0) {
3     printf("%d\n", p->f->x);
4   }
5 }
```

- $\rightarrow$ *conservatively* estimate objects that *may be* accessed in future
  - NEVER reclaim those that are accessed
  - OK not to reclaim those that are in fact never accessed
- in the above example, OK to retain objects pointed to by `p` when the line 2 is about to start

Objects that “may be” accessed

- global variables
- local variables of active function calls (calls that have started but have not finished)

```
int * s, * t;
void h() { ... }
void g() {
  ...
  h();
  ... = p->x ...  }
void f() {
  ...
  g()
  ... = q->y ... }
int main() {
  ...
  f()
  ... = r->z ... }
```



Objects that “may be” accessed

- global variables
- local variables of active function calls (calls that have started but have not finished)
- objects reachable from them by traversing pointers

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int * s, * t;
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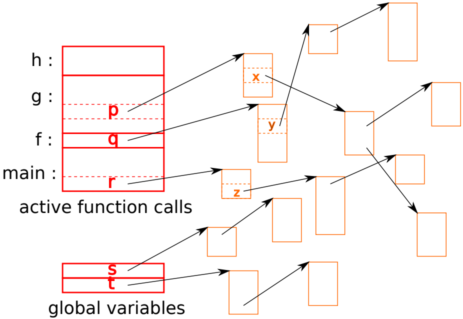

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- an object: the unit of automatic memory allocation/release (malloc in C; objects in Java; etc.)

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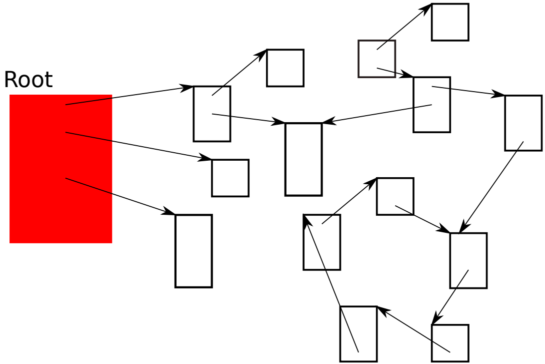
> the basic principle of GC:
> objects unreachable from the root are dead

The two major GC methods

- traversing GC:
  - simply traverse pointers from the root, to find (or *visit*) objects reachable from the root
  - reclaim objects not visited
  - two basic traversing methods
    - mark&sweep GC
    - copying GC
- reference counting GC (or RC):
  - during execution, maintain the number of pointers (reference count) pointing to each object
  - reclaim an object when its reference count drops to zero
  - note: an object's reference count is zero → it's unreachable from the root
- remark: "GC" sometimes narrowly refers to traversing GC

How traversing GC works

- traverse pointers from the root
- once all pointers have been traversed, objects that have not been visited are garbage
- the difference between mark&sweep and copying is covered later



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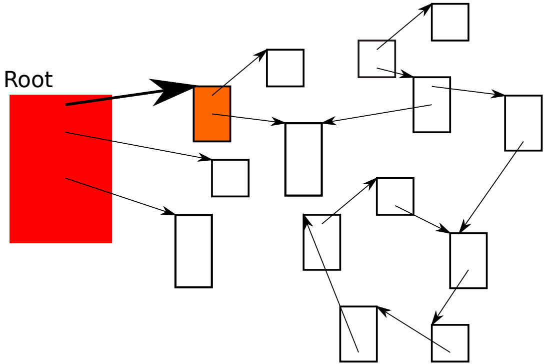

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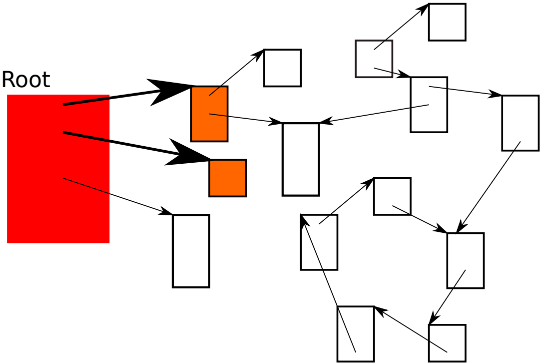

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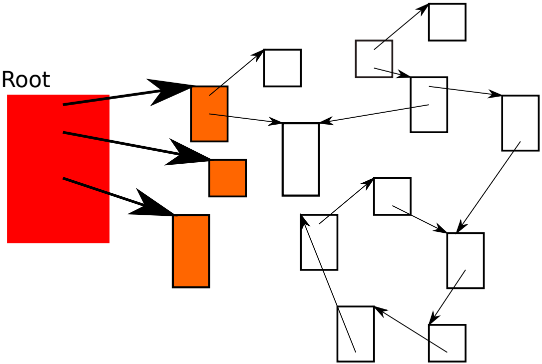

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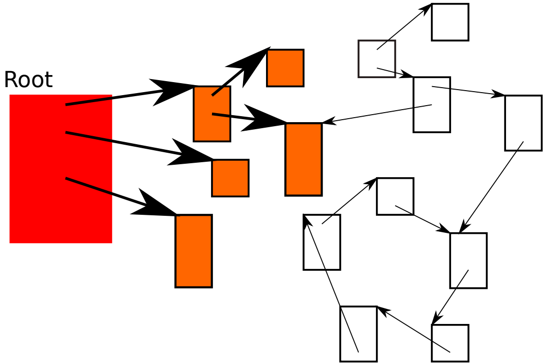

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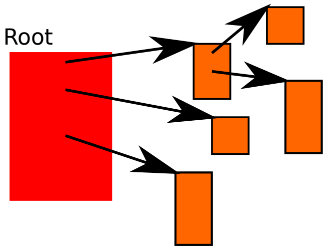
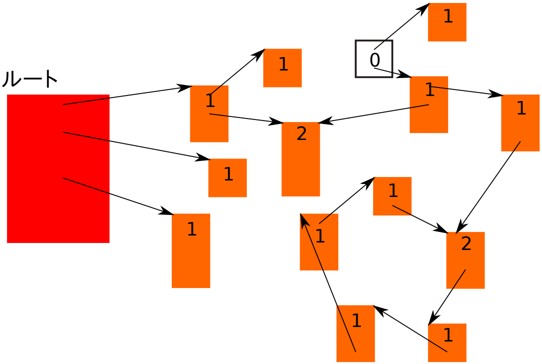


How reference counting works

- each object has a reference count (RC)
- update RCs during execution; e.g., upon `p = q;` →
  - the RC of the object `p` points to `-=` 1
  - the RC of the object `q` points to `+=` 1
- reclaim an object when its RC drops to zero → RCs of objects pointed to by the now reclaimed object decrease



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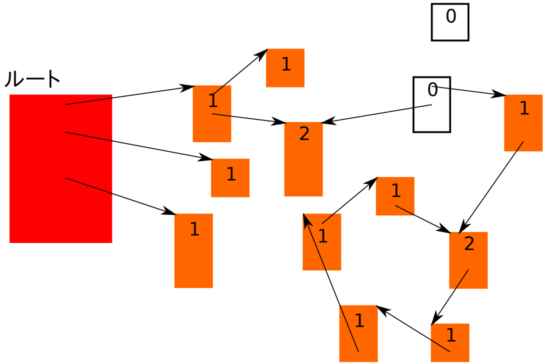

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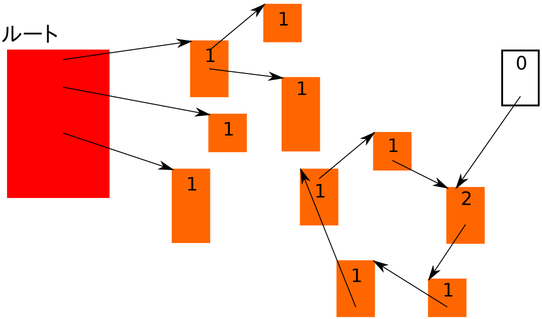

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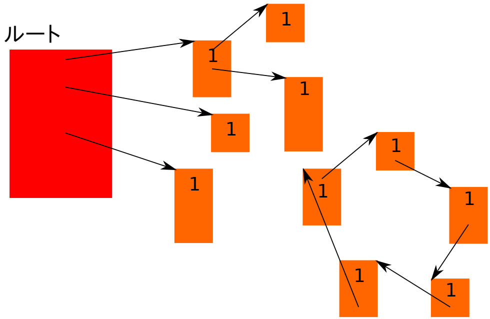

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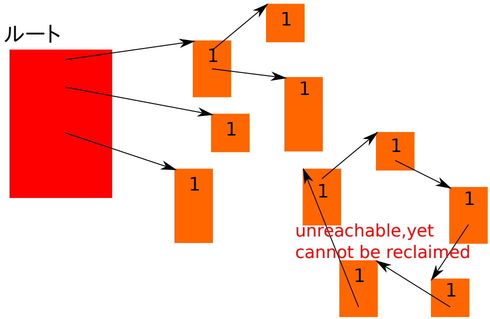

When an RC changes

- a pointer is updated `p = q`; `p->f = q`; etc.
- a function gets called

```
int main() {
  object * q = ...;
  f(q);
}
```

- a variable goes out of scope or a function returns

```
f(object * p) {
  ...
  {
    object * r = ...;

  } /* RC of r should decrease */
  ...
  return ...; /* RC of p should decrease */
}
```

- etc. any point pointer variables get copied / become no longer used

GC will be covered more deeply in later weeks